

THE COULATZ CONJECTURE STATES THAT IF YOU PICK A NUMBER, AND IF ITSEVEN DIVIDE ITBY TWO AND IF IT'S OOD MULTIPLY IT BY THREE AND ACD ONE AND YOO REPEAT THIS PROCEDURE LONG ENOUGH, EVENTOALLY YOUR FRIENDS WIL STOP CACUING TO SEE 脤 YOU WANT TO HANG OUT.

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## Ways of Presenting Algorithms (Algorithms Part II)

ver. 14 z drobnymi modyfikacjami!

## Wojciech Myszka

2023-11-21 07:29:41 + 0100


## Expressing/Presenting algorithms

1. Natural language. Use simple sentences in the imperative mood.

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2. Flowcharts. In a moment...

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3. Decision tables. A little bit later. . .

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## Expressing/Presenting algorithms

1. Natural language. Use simple sentences in the imperative mood.
2. Flowcharts. In a moment...
3. Decision tables. A little bit later...
4. Pseudocode - formal type of writing, similar to...
5. Programing languages. Never. . . ?! (But remember about Blockly...)

## Algorithm

## Algorithm <br> a step-by-step procedure for solving a problem or accomplishing some end

## Expressing algorithms

## Flowchart


#### Abstract

Flowchart (also block diagram) A flowchart is a type of diagram that represents an algorithm or process, showing the steps as boxes of various kinds, and their order by connecting these with arrows. This diagrammatic representation can give a step-by-step solution to a given problem. Process operations are represented in these boxes, and arrows connecting them represent the flow of control.


Block diagram reveals important steps in the algorithm and the logical relationships between them.

Introduced in the early twenties of the XX century. Used, among others, by von Neumann.

## Symbols used in the flowcharts

Common shapes


## Flowchart template



## Flowchart

Terminator Symbol

Terminator symbol signals the start or end of a process. It usually contains the word "Start", "Begin" or "Stop", "End".

## Flowchart

Data symbol

Data Symbol (also Input/Output) The Data flowchart shape indicates inputs to and

## Data Symbol

 outputs from a process. As such, the shape is more often referred to as an I/O shape than a Data shape. READ z, WRITE z+10.
## Flowchart

## Process Symbol Generic

 processing step. Examples: "Add 1 to $z$ " or " $\mathrm{z}=\mathrm{z}+\mathbf{1}^{\prime \prime}$; "replace identified part"; "save changes" or similar.
## Flowchart

Predefined process
(Subroutine) is used to show complex processing steps which
 may be detailed in a separate flowchart.

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## Flowchart

Decision Symbol shows where a decision is necessary, commonly a Yes/No question or True/False test. For example: Check if $a$ is
 equal to $b$, if YES do something if NO do something else.

## Flowchart

Connector If you need to connect to another page or another section of the chart, and can't draw a line, you can use a circle. You draw the line to the circle and label the circle with a letter or other symbol: 4.3, 2, B2.

## Selecting the highest person in the room

Natural language

1. Select any person from the audience, think of it as the highest (and place it near the door).

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2. Are there any people left in the room? if so - go to step 5 .
3. If not - the highest person is one standing near the door.

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4. End of the algorithm

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4. End of the algorithm
5. Take another person from the room.

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4. End of the algorithm
5. Take another person from the room.
6. Compare it with one standing near the door - is it higher?

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4. End of the algorithm
5. Take another person from the room.
6. Compare it with one standing near the door - is it higher?
7. If not go to step 2

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1. Select any person from the audience, think of it as the highest (and place it near the door).
2. Are there any people left in the room? if so - go to step 5 .
3. If not - the highest person is one standing near the door.
4. End of the algorithm
5. Take another person from the room.
6. Compare it with one standing near the door - is it higher?
7. If not go to step 2
8. If so - replace the person standing near the door, go to step 2

## Selecting the highest person in the room

Flowchart

Start

## Selecting the highest person in the room

Flowchart


## Selecting the highest person in the room

Flowchart


## Selecting the highest person in the room

Flowchart


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## Selecting the highest person in the room

## Flowchart



## Selecting the highest person in the room

## Flowchart



## Salary



1. make a note of the number 0;
2. proceed through the list, adding each employee's salary to the noted number;
3. having reached the end of the list, produce the noted number as output.

## Euclidean Algorithm

Natural Language

1. [Find the reminder] Divide $m$ by $n$ let $r$ be the reminder. (We have $0 \leq r<n$.)
2. [Is zero?] If $r=0$ finish the procedure; the answer is $n$.
3. [Simplifying] Let $m \leftarrow n, n \leftarrow r$ and came back to the step 1 .

## Euclidean Algorithm

Flowchart


## Decision table

- An alternative to the flowchart.
- Tables are a quick and easy way for humans to read, understand and execute a complex procedure.
- Work best with complex decision problems.
- It is not suitable to describe computational problems.
- Today it is a little forgotten.


## Decision table cont.

The idea comes from the '50s (20th century!)

## Structure of DT

| Conditions | Condition alternatives |
| :--- | :--- |
| Actions | Action entries |

- Each decision corresponds to a variable whose possible values are listed among the condition alternatives.
- Each action is a procedure or operation to perform, and the entries specify whether (or in what order) the action is to be performed for the set of condition alternatives the entry corresponds to.


## Example

|  |  | Rules |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Printer does not print | Y | Y | Y | Y | N | N | N | N |
|  | A red light is flashing | Y | Y | N | N | Y | Y | N | N |
|  | Printer is unrecognized | Y | N | Y | N | Y | N | Y | N |
| $\begin{array}{\|l\|l} \hline \frac{n}{\partial} \\ \stackrel{\rightharpoonup}{4} \\ \hline \end{array}$ | Check the power cable |  |  | X |  |  |  |  |  |
|  | Check the printer-computer cable | X |  | X |  |  |  |  |  |
|  | Ensure printer software is installed | X |  | X |  | X |  | X |  |
|  | Check/replace ink | X | X |  |  | X | X |  |  |
|  | Check for paper jam |  | X |  | X |  |  |  |  |

## Buing a car

|  | an excessive price | price OK | underestimated price |
| :--- | :--- | :--- | :--- |
| The car meets all require- <br> ments |  |  |  |
| The car does not meet all <br> the requirements but the <br> condition and equipment <br> are acceptable |  |  |  |
| The car does not meet the <br> requirements |  |  |  |

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| The car does not meet all <br> the requirements but the <br> condition and equipment <br> are acceptable |  |  |  |
| The car does not meet the <br> requirements |  |  |  |

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| The car does not meet all <br> the requirements but the <br> condition and equipment <br> are acceptable | Resign |  |  |
| The car does not meet the <br> requirements |  |  |  |

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| The car does not meet all <br> the requirements but the <br> condition and equipment <br> are acceptable | Resign | Bargain; Buy if the price is <br> reduced |  |
| The car does not meet the <br> requirements |  |  |  |

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## Buing a car

|  | an excessive price | price OK | underestimated price |
| :--- | :--- | :--- | :--- |
| The car meets all require- <br> ments | Bargain; if you do not suc- <br> ceed come back the next <br> day and bargain again; <br> buy even if you can not re- <br> duce the price | Bargain; Buy It whatever <br> the result of negotiations | Buy |
| The car does not meet all <br> the requirements but the <br> condition and equipment <br> are acceptable | Resign | Bargain; Buy if the price is <br> reduced | Bargain; Buy regardless of <br> the results the negotia- <br> tions |
| The car does not meet the <br> requirements |  |  |  |

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## Buing a car

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| The car does not meet the <br> requirements | Resign |  |  |

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| The car does not meet the <br> requirements | Resign | Resign |  |

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| The car does not meet all <br> the requirements but the <br> condition and equipment <br> are acceptable | Resign | Bargain; Buy if the price is <br> reduced | Bargain; Buy regardless of <br> the results the negotia- <br> tions |
| The car does not meet the <br> requirements | Resign | Resign | Bargain the extras; Buy if <br> the price of the extras is <br> reasonable |

## Dressing a coat

|  |  | R1 | R2 | R3 |
| :--- | :--- | :---: | :---: | :---: |
| C1 | Raining | Y | Y | N |
| C2 | Cold | Y | N | Y |
| A1 | Dress warm raincoat | X |  |  |
| A2 | Dress raincoat |  | X |  |
| A3 | Dress warm coat |  |  | X |

## Dressing the coat cont.

In the table we have omitted the condition:
Raining N
Cold N
it does not require any special action, although it could be added by defining the action: 'Do not wear any coat'.

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## Internet shop (B2B)

Another example is the decision table describing the activities related to the adoption and execution of the contract.

The table includes also the shop policy that can be described as follows:

1. The company serves registered customers only.
2. The company supplies goods from the list of goods only.

## Internet schop cont.

| C1 | Item on the list? | Y | N | - | Y |
| :--- | :--- | :---: | :---: | :---: | :---: |
| C2 | Customer registered? | Y | - | N | Y |
| C3 | Sufficient supply of goods? | Y | - | - | N |
| A1 | Book the item | X |  |  |  |
| A2 | Log the transaction | X |  |  |  |
| A3 | Save the item for the list to realize |  |  |  | X |
| A4 | Deliver the item | X |  |  |  |
| A5 | Reject the transaction |  | X | X |  |

"-" means Do not care

## Programming Languages

Flowchart


## Programming languages

Blockly


## Programming languages

Java Script

```
var m, n, r;
m = parseFloat (window. prompt('input m'));
n = parseFloat(window.prompt('input n'));
r = m % n;
while (r != 0) {
        m = n;
        n = r;
        r = m % n;
}
window.alert(n);
```


## Programming languages

## Python

$$
\begin{aligned}
& \mathrm{m}=\text { None } \\
& \mathrm{n}=\text { None } \\
& r=\text { None } \\
& m=\text { float (text_prompt('input_m')) } \\
& \mathrm{n}=\mathrm{float}(\text { text_prompt('input」n')) } \\
& r=m \% n \\
& \text { while } r \text { ! }=0 \text { : } \\
& \mathrm{m}=\mathrm{n} \\
& \mathrm{n}=\mathrm{r} \\
& r=m \% n \\
& \text { print(n) }
\end{aligned}
$$

## Programming languages

PHP

```
$m;
$n;
$r;
$m = floatval(readline('input mm'));
$n = floatval(readline('inputьn'));
$r = $m % $n;
while ($r != 0) {
    $m=$n;
    $n=$r;
    $r=$m % $n;
}
print($n);
```


## Programming languages

Lua

$$
\begin{aligned}
& \mathrm{m}=\text { tonumber(text_prompt('input m'), 10) } \\
& \left.\left.\mathrm{n}=\text { tonumber(text_prompt('input } \mathrm{n}^{\prime}\right), 10\right) \\
& \mathrm{r}=\mathrm{m} \% \mathrm{n} \\
& \text { while } r^{\sim}=0 \text { do } \\
& \quad \mathrm{m}=\mathrm{n} \\
& \mathrm{n}=\mathrm{r} \\
& \quad \mathrm{r}=\mathrm{m} \% \mathrm{n} \\
& \text { end } \\
& \text { print }(\mathrm{n})
\end{aligned}
$$

## Programming languages

## Dart

```
var m, n, r;
main() {
m = Math.parseDouble(Html.window.prompt('input m', ''));
n = Math.parseDouble(Html.window.prompt('input n', ''));
r = m % n;
while (r != 0) {
        m = n;
        n = r;
        r = m % n;
    }
print(n);
```


## Control structure

- An algorithm can be thought of as being executed by a little robot, or a processor). The processor receives orders to run around doing this and that, where the "this and that" are the basic actions of the algorithm.
- It should be quite obvious that the order in which the basic actions are carried out is crucial.
- Creating an algorithm we will use the sequence of commands. We assume that after one command will be executed the next.
- Sometimes we will like to change the execution order of commands using conditionals if something (happens) do this else do that

Direct sequencing
Direct sequencing
of the form "do A followed by B," or "do A and then B." (Every semicolon or period in the recipe hides an implicit "and then" phrase, for example, "gently fold in chocolate; [and then] reheat slightly . . .")

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| move forward |
| :--- |
| turn left $\cup \vee$ |
| move forward |
| turn right $\circlearrowright ~ \nabla$ |
| move forward |

## Conditional branching

## Conditional branching

of the form "if $Q$ then do $A$ otherwise do $B$," or just "if $Q$ then do $A$," where $Q$ is some condition. (For example, in the recipe "reheat slightly to melt chocolate, if necessary," or "serve with whipped cream, if desired.")

## Conditional branching

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## |terations |

Sequencing and branching, do not explain how an algorithm of fixed-maybe even short-length can describe processes that can grow increasingly long, depending on the particular input.

## Bounded iteration

of the general form "do A exactly N times," where N is a number.

## Conditional iteration

sometimes called unbounded iteration, of the form "repeat A until Q," or "while $Q$ do $A$," where $Q$ is a condition. (For example, in the recipe "beat egg whites until foamy.")
|terations can be nested!

## Bounded iteration

## Factorial

```
set n prompt for number with message Get n
set S , 1
count with i from 1 to get n
do set s get s }
set z n!=
to z}\mathrm{ append text , get s
print get z
```

Iteration cont.
Nested iterations


Iteration cont.
Nested iterations


## GoTo Instruction

## GoTo

Instruction has the general form "goto G," where G marks some point in the text of the algorithm.

Many researchers are opposed to using "goto"s freely in algorithms.

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## Subroutines



## Subroutines


draw a circle
(ㄷ) (3) to draw a circle
repeat 36 times
do move forward v by 10


## Subroutines I

1. We are given a lengthy text and we are interested in finding out how avaricious its author is by counting the number of sentences that contain the word "money."
2. We are not interested in the number of times the word "money" occurs, but in the number of sentences in which it occurs.
3. An algorithm can be designed to run through the text looking for "money."
4. Upon finding such an occurrence, it proceeds to run ahead looking for the end of a sentence, which for our purposes is assumed to be a period followed by a space; that is, the ". "
5. When the end of a sentence is found, the algorithm adds 1 to a counter which was initialized to 0 at the start.
6. Algorithm then resumes its search for "money" from the beginning of the next sentence.

## Subroutines



## Subroutines



## Subroutines



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## Subroutine

## Subroutine

(also termed procedure, function, routine, method, or subprogram), is a part of source code (algorithm) within a larger computer program that performs a specific task and is relatively independent of the remaining code (algorithm).

## Variables I

- A variable is not a number, a word, or some other item of data. Rather, it can be treated as a small box, or cell, or place, in which a single item can be kept.
- Algorithms utilize variables (with different names) for different purposes, bat generally for storing some values. They are kind of a storage or a memory.
- New value putted into the variable erases its previous content.
- Taking a value from the variable does not change its content.
- Each variable can contain only one value at a time.
- Notation similar to mathematical is used to describe operations on variables: $X=5, X:=X+1$, or $X \leftarrow A+B$


## Vector, array I

- Let us think about our employee list. It may be viewed simply as a multitude of data elements, which we might decide to keep, or store, in a multitude of variables, say $X, Y, Z, \ldots$
- Storing all salaries that way is not very suitable for the mentioned algorithm: each element in the list would have to be referred to in the algorithm by a unique name.
- We need lists of variables that can be "run through," or accessed in some other way, but without the need to name each of their elements explicitly.
- In mathematics, such elements are called "vectors", or one-dimensional arrays.
- The name is assigned to all structures (containing all data).


## Vector, array II

- We can access each structure element, for example, by its position on the list, we note this: $V(I)$ means the content of the $I$-th element of the vector $V$. The alternative notation looks like this: $V[J]$. $I, J$, are so-called pointers or indexes

Here each cells represents one elementary variable.

| $V(1)$ | $V(2)$ | $V(3)$ | $V(4)$ | $V(5)$ |
| :--- | :--- | :--- | :--- | :--- |

In mathematical notation:

| $V_{1}$ | $V_{2}$ | $V_{3}$ | $V_{4}$ | $V_{5}$ |
| :--- | :--- | :--- | :--- | :--- |

## Simple algorithm using vectors

1. do the following $N-1$ times:
$1.1 X \leftarrow 1$;
1.2 while $X<N$ do the following:
1.2.1 if $V[X+1]<V[X]$ then exchange them;
1.2.2 $X \leftarrow X+1$.

## Simple algorithm using vectors

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1.2.2 $X \leftarrow X+1$.

There are also special "indexed" versions of iterative control constructs, tailored towards vector traversal. For example, we can write:
for $X$ going from 1 to 100 do the following
which is similar to:
do the following 100 times

## Array (multidimensional) I

- In many cases, it is convenient to arrange the data, not in a simple, one-dimensional list, but in a table.
- The corresponding algorithmic data structure is called a matrix, or a two-dimensional array, or simply an array for short.
- The standard second-grade multiplication table is a 10 by 10 array in which the data item at each point is the product of the row and column indices;
- Referring to an array element is typically achieved using two indices, row and column. We write $A[5,3]$ for the element located in row 5 and column 3. Sometimes such referring will be noted: $A(5,3)$ or $A[5][3]$.


## Array (multidimensional) II

- Running through the entire array can be achieved by an outer loop running through all rows and an inner one running through all of a particular row elements, or vice versa.

Mathematical notation:

$$
\begin{array}{lll}
v_{1,1} & v_{1,2} & v_{1,3} \\
v_{2,1} & v_{2,2} & v_{2,3} \\
v_{3,1} & v_{3,2} & v_{3,3} \\
v_{4,1} & v_{4,2} & v_{4,3}
\end{array}
$$

Algorithmic notation:

$$
\begin{array}{lll}
V(1,1) & V(1,2) & V(1,3) \\
V(2,1) & V(2,2) & V(2,3) \\
V(3,1) & V(3,2) & V(3,3) \\
V(4,1) & V(4,2) & V(4,3)
\end{array}
$$

## Control structures for arrays



## Queue

- An interesting variation on the vector/array.
- Sometimes a list is used just to model a queue, in which case all the algorithm needs in way of interaction with the list are the ability to add elements to its "back" and remove them from its "front."
- The structure is called FIFO (First In First Out)



## Stack

- A structure similar to a vector (list)
- Adding elements to "front"
- Remove also from "front".



## Tree or hierarchy

One of the most important and prominent data structures in existence is the tree.

- Structure with the order.
- There is a special object that is the "beginning" of the whole structure: root.
- Other elements are called descendants or offspring.
- Each object can have a number of equivalent descendants.
- Each offspring is called a node.
- Nodes at the "end" of the tree, having no offspring are leaves.
- Branch sequences of nodes corresponding to downward traversals in the direction from the root to a leaf.


## Tree



## Other data structures

- lists (some similarity to vectors or arrays).
- databases (some similarity to arrays).
- graphs (some resemblance to trees).


## Part I

## Excursus

## Selecting the highest person

(in the room)

1. Height of all people are saved in an array called $\mathbf{H}$. (This should be height, but it takes up too much space.)
2. Variable $\mathbf{N}$ contains the number of persons in the room (the length of array H).
3. Variable MAX contains the maximum value from array $H$ (when the algorithm ends) i.e. height of the highest person in the room.
4. I - auxiliary variable.

Maximum value in an array

Maximum value in an array

MAX
$\mathrm{I}=1$

Maximum value in an array


## Maximum value in an array



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This works for odd positive values and does not work for even values.

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Let's start from 7: 7, 22, 11, 34, 17, 52, 26, 13, 40, 20, 10, 5, 16, 8, 4, 2, 1 And what with other values?

## A little bit more complicated problem... I

Collatz conjecture

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## A little bit more complicated problem... II

## Collatz conjecture

It turns out that the above algorithm either completes relatively quickly or it generates a very long (an infinite?) chaotic sequence of numbers.

No one was able to prove that sequence generated by the algorithm begin to repeat (which means the algorithm will not stop never) or prove that for some particular initial value $X$ the algorithm will stop.

## A little bit more complicated problem... III

## Collatz conjecture

The Collatz conjecture is: This process will eventually reach the number 1 , regardless of which positive integer is chosen initially.

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