

Floating-point numbers ver. 11 z drobnymi modyfikacjami!

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Main milestones in building modern computers:

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Problems:

- limited number of bits for storing values
- ▶ 8, 16, 32, 64, 128,...
- when the result of arithmetic operation does not "fit" result is wrong! This is called *overflow*.





How to store non-integer numbers?

In spreadsheet we can choose between:

- ▶ general: 1.5782
- number 1.58
- percent 157.82%
- ▶ currency 1.58 €
- scientific 1.58E+00
- thousand separator

This is only external representation! What about internal representation?



Fixed point

Let's assume 16 bits non-integer arithmetic.

- 1. First eight bits for the integer part
- 2. Second eight bits for the fractional part

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
sign			inte	ger	part					fra	action	nal pa	art		



Floating-point

All sixteen bits (in general 32 or 64) are used for storing

- 1. sign
- 2. fractional part x
- 3. exponent e

and the number is represented as $\mathit{sign} \cdot x \cdot 2^e$

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	
									,	 						

sign (e) exponent (x) fractional part



Fixed-point vs Floating-point

- Fixed point means that we are using a fixed number of digits for remembering the integer part of a number and a fixed number for the non-integer (fractional) part, for example, 6 (integer) + 2 (fractional) in "financial calculator" (for home use). So the position of the decimal point is also fixed.
- Floating-point means that position of the decimal point changes during calculations: we have a fixed number of significant digits.



Fractions

In decimal numbers we have: $345.5=3*10^2+4*10^1+5*10^0+5*10^{-1}$ So, analogically we can write: $101011001.1_{(2)}=2^8+2^6+2^4+2^3+2^0+2^{-1}$

Let's think about 3rd (binary) digit (in fractional part) $0.001_{(2)} = 0.125_{(10)}$ To store fractions with reasonable precision this needs a lot of bits. The same when there is a lot of digits on the left of "digital point".



Big decimal numbers

To store big number one can use "scientific notation" (or floating-point). Instead of writing

c = 299792458 m/s

one can write	
	$c = 2.99792458 * 10^8 \mathrm{~m/s}$
or	a - 2 00702459 E 8 m /a
or (approximately)	$c = 2.99792458 \to 8 \text{ m/s}$
or (approximately)	$c=3 \mathrm{E}8 \mathrm{~m/s}$
2.99792458 is sometimes call	led a significand (mantissa); 8 is called an exponent.

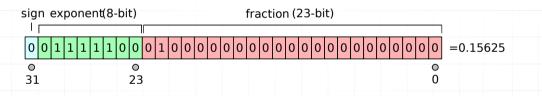


"Big" binary digits

One can use the same way

$$1.11_{(2)} * 2^7$$

Floating-point numbers are described in IEEE-754 Standard (established in 1985):



All floating-point numbers are stored in, so-called, *normalized form*, i.e. there is only one digit on the left side of the decimal point, and is greater than zero.



Precision

32 bits numbers

- "Binary" precision: 24 bits
- ▶ "Decimal" precision: ≈ 7.2 decimal digits (this means: mostly 7, sometimes 8, in average 7.2).

64 bits numbers

- "Binary" precision: 53 bits
- "Decimal" precision: \approx 15.9 decimal digits.



Range

- $1.\,$ 32-bits: $1.17549 \:\mathrm{E}-038$ to $3.40282 \:\mathrm{E}+038$
- 2. 64-bits: 2.22507 $\mathrm{E}-308$ to 1.79769 $\mathrm{E}+308$



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Floating-point values...

... are **rational numbers**. They can be expressed as the quotient or fraction $\frac{p}{q}$ of two integers, a numerator p and a non-zero denominator q.



Floating-point arithmetic I

Let's assume that our computer uses decimal floating-point arithmetic with three digits.

1. Multiplication.

Easy: multiplying of mantissas and adding exponents

 $1.33 \, e + 3 * 1.55 \, e + 7 = 2.0615 \, e + 10$

Next, the result has to be "shorten" (cut) to three digits: 2.06 e+10

Note: We are loosing value of $0.0015 e+10 = 15\,000\,000$ (Yes! Fifteen millions!) Caution: Sometimes strange happens. After the operation result is **denormalised**: more then one figure on the left to the digital point. Result is normalized

(exponent is corrected), and rounded :

5.55 e+0 * 6.33e+0 = 35.13 e+0 = 3.51 e+1

2. Division:

Like multiplication: mantissas divided, exponent subtracted.

1.33 e + 0/9.88 e + 0 = 0.134615385 e + 0 = 1.35 e - 1

(result normalized and rounded).



Floating-point arithmetic II

3. Addition.

A simple method to add floating-point numbers is to first represent them with the same exponent (denormalize!) 1.22 e+0 + 3.35 e - 4 = 1.22 e+0 + 0.000335 e+0 = 1.220335 e+0 = 1.22 e+0and next normalize and round... But check carefully the result!

- 4. Subtraction.
 - Like addition.



Some problems...

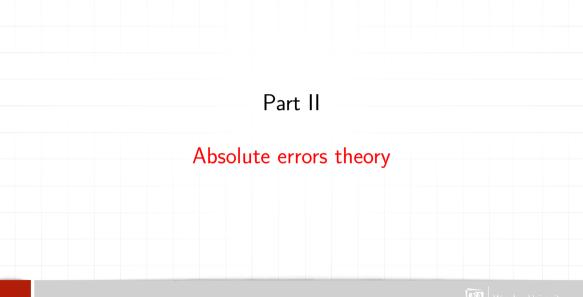
- 1. The limited number of bits used to store numbers! (There are special applications allowing for arithmetic with arbitrary numbers of digits.)
- "Overflow" occurs when the result of arithmetic operation does not "fit" in a word (32 or 64 bits).
- 3. Most of the numbers that (out of habit), are considered to be accurate, do not have exact binary representation (0.5 is OK but 0.1 NO).



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- 4. The last one is OK. We do remember fraction $rac{1}{3}=0.3(ar{3})$, but...





Wrocław University of Science and Techno

Basic definitions

Quantity

any mathematical constant, the result of some mathematical operations (actions), a root of the equation solved. π is defined as the ratio of the circumference of a circle to its diameter, $\sqrt{2}$ is a root of the quadratic equation $x^2 - 2 = 0$ (or the diagonal of a unit square).





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value derived directly from the definition, not burdened by any errors.



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The approximate value of the quantity

numerical value obtained by calculation. Typically, the calculation did not get the exact value.

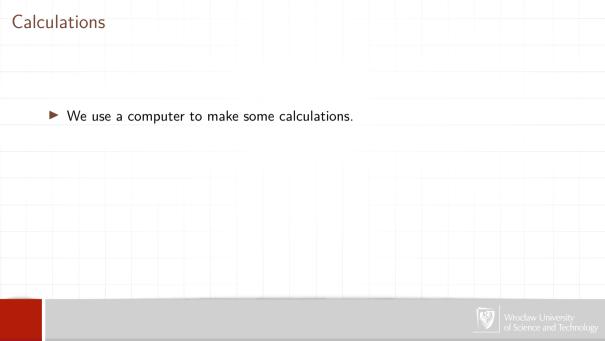


Pressure, temperature, length, concentration — are examples of physical quantities, which often are measured.

Each measurement was burdened by an **error** resulting from the **accuracy** of the measurement tool used.

- So, for example:
 - Quantity: the temperature at some point in the room.
 - ▶ The exact value of the quantity: the temperature at this particular point
 - The approximate value of the quantity: temperature measured with a thermometer.





Calculations

- ▶ We use a computer to make some calculations.
- The computer gives the result (a = 5.34273343) with 8 digits after the decimal point.



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Calculations

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- ▶ Can we say that the result has all the figures **correct**? I mean that the difference between the result and the exact value Is less than 0.5×10^{-8} ?
- What about the situation that the method of calculation is inaccurate?



The equation describing the period of oscillations (in seconds) looks like:

$$\overline{f} = 2\pi \sqrt{\frac{L}{g}}$$
 (1)

where L is the length (in meters) and g is a gravitational acceleration constant.



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$$\Gamma = 2\pi \sqrt{\frac{L}{g}} \tag{1}$$

where *L* is the length (in meters) and *g* is a gravitational acceleration constant. What is the value of π ? What is the value of *g*?



The equation describing the period of oscillations (in seconds) looks like:

$$T = 2\pi \sqrt{\frac{L}{g}} \tag{1}$$

where L is the length (in meters) and g is a gravitational acceleration constant. What is the value of π ? What is the value of g? Let's assume, that L = 1 m.



Mathematical pendulum

	Т	pi	g	L
	.00504969	3.14	9.81	1
	.00632679	3.142	9.81	1
	.00607137	3.1416	9.81	1
	.00606499	3.14159	9.81	1
	.00606690	3.141593	9.81	1
	.00606668	41592654	9.81	1
	.00640929	41592654	9.80665	1
	.00368655	41592654	9.83332	1
е	.00911030	41592654	9.7803	1
	.00591333	41592654	9.8115	1
	.00583156	41592654	9.8123	1
	.00601556	41592654	9.8105	1
arct	.00364783	41592654	9.8337	1





The absolute error of the approximate (measured) value of the quantity I

Let A be the exact value of the quantity, and a be its approximate value.

Absolute error

is any number Δa satisfying the condition:

$$|A-a| \leq \Delta a$$
,

i.e. such a number, that:

$$a-\Delta a\leq A\leq a+\Delta a.$$

Approximate value *a* and its absolute error Δa determine an interval:

$$<$$
 a Δ a; a $+$ Δ a $>$

to which belongs the exact value A.

he absolute error is not specified unambiguously!



Rough value

Rough value

If a is an approximate value of quantity A with an error Δa then

∆a a

I will call **rough value** for A.



We know that $\pi = 3.14159265...$ 3.14 is commonly used as an approximation of π in calculations. What is an absolute error of this approximation?





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Because

 $3.14 - 0.0016 \le \pi \le 3.14 + 0.0016$

or

 $3.1384 \leq \pi \leq 3.1416$

the absolute error of this approximation is 0.0016. Exact value of π is somewhere in the interval

< 3.1384; 3.1416 >

so we can write $\pi = 3.14^{0.0016}$



"Approximate equality"

If two rough values $\stackrel{\alpha}{a}$ and $\stackrel{\beta}{b}$ are such, that the interval $\langle a - \alpha; a + \alpha \rangle$ is **included** in the interval $\langle b - \beta; b + \beta \rangle$ we can tell that $\stackrel{\alpha}{a}$ is approximately equal to $\stackrel{\beta}{b}$. We will note this as: $\stackrel{\alpha}{a} \stackrel{\beta}{\Rightarrow} \stackrel{\beta}{h}$

The fact that $\stackrel{\alpha}{a}$ is approximately equal to $\stackrel{\beta}{b}$ does not imply the opposite: that $\stackrel{\beta}{b}$ is approximately equal to $\stackrel{\alpha}{a}!$ (The relation is not reflexive.)



For any rough value $\stackrel{\alpha}{a}$ and any real value b following relation is fulfilled:

 $\stackrel{lpha}{a} \stackrel{lpha+|a-b|}{b}$

this means that $\stackrel{lpha}{=}$ is approximately equal to $\stackrel{lpha+|a-b|}{b}$



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So, instead of using 3.14159265 as the π I can use simply 3 but...



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... I should tell that the error is |3.14159265 - 3| + 0.00000003 = 0.141592653!

Calculations are simpler but burdened with 47197551 times greater absolute error.



Rounding rough values cont.

Of course, the following rounding is "better":

We are rounding numbers for practical reasons: when the result has too much significant digits. . .



Rounding rough values cont.

When a = b and $\beta \ge \alpha$, we can write that:

 $\stackrel{\alpha}{a} \Rightarrow \stackrel{\beta}{b}$

So $\begin{array}{c} 0.0015927 & 0.0016 \\ 3.14 \Rightarrow 3.14 \end{array}$

These values are approximately equal.



Rounding rules I

- When the result of calculations has a lot of digits we can remove (cut off) the least significant digits — increasing rounding error.
- When the first of removed digits is 0, 1, 2, 3, 4 the last remaining digit is not changed.
- ▶ When the first of removed digits is 5, 6, 7, 8, 9 we are increasing the result by one on the least significant digit.

These rules are sometimes called "proper rounding," but easily one can find other rules:

- Round half up
- Round half down
- Round half away from zero
- Round half towards zero
- Round half to even (banker's rounding)

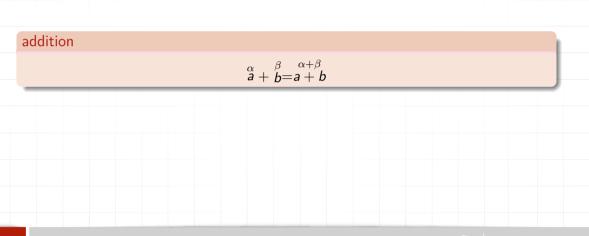


Rounding rules II

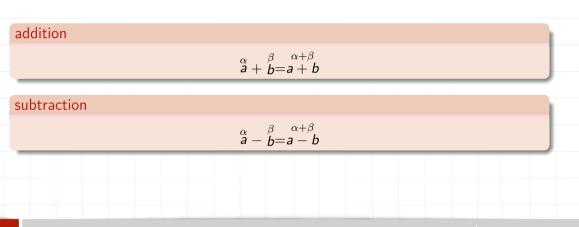
- Round half to odd
- Stochastic rounding

Homework: Read about rounding rules and decide which one (of the above mentioned rules) is "proper rounding?"

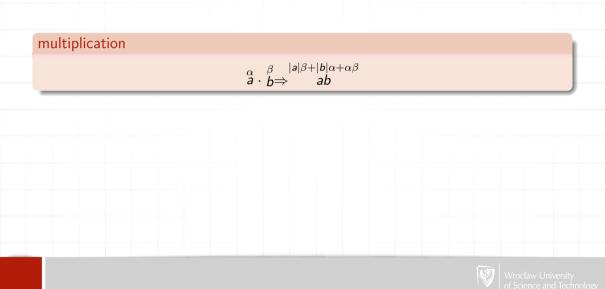


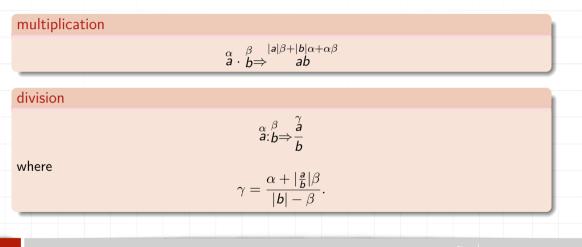














Error in other mathematical operations:

- 1. rising to power?
- 2. square root?
- 3. trigonometric functions?
- 4. other...?



Arithmetic operations on rough values addition

Do we need to memorize all these rules? Better is to understand them.

1. The first "worst case":

$$a - \alpha + b - \beta = (a + b) - (\alpha + \beta)$$

2. The second "worst case":

$$\mathbf{a} + \alpha + \mathbf{b} + \beta = (\mathbf{a} + \mathbf{b}) + (\alpha + \beta)$$



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(Homework: How it will work for subtraction? Multiplication? Division?)



Example

Calculate the value of the polynomial

$$w(x) = a_0 x^4 + a_1 x^3 + a_2 x^2 + a_3 x + a_4$$

for x = 2.1.

Let's assume, those polynomial coefficients are exact values, and are equal to:

$$a_0 = 2.3, a_1 = 3, a_2 = -4.5, a_3 = 7.2, a_4 = -0.1$$

First, we perform calculations accurate to two decimal places, and then to four.



Example, cont. I

two decimal places

$$x^{2} = 2.1 \times 2.1 = 4.41$$

$$x^{3} = 4.41 \times 2.1 = 9.261 \Rightarrow 9.26$$

$$x^{4} = 9.26 \times 2.1 \Rightarrow 19.446 \Rightarrow 19.45$$

$$2.3 \times x^{4} = 2.3 \times 19.45 \Rightarrow 44.735 \Rightarrow 44.74 \Rightarrow 44.74$$

$$3x^{3} = 3 \times 9.26 \Rightarrow 27.78$$

$$-4.5x^{2} = -4.5 \times 4.41 \Rightarrow -19.845 \Rightarrow -19.85$$

$$7.2x = 7.2 \times 2.1 \Rightarrow 15.12$$

sum:

w(2.1) = 44.74 + 27.78 - 19.85 + 15.12 - 0.1 = 67.69

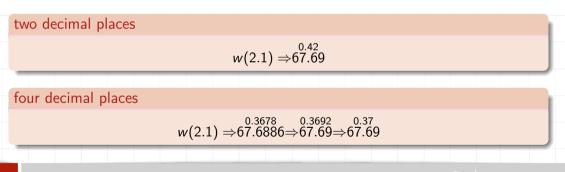


Example, cont. I four decimal places $x^2 = 21 \times 21 = 441$ $x^{3} = 4.41 \times 2.1 = 9.261$ $x^4 = 9.261 \times 2.1 = 19.4481$ $2.3 \times x^4 = 2.3 \times 19.4481 = 44.73063 \Rightarrow 44.7306$ $3x^3 = \overset{0}{3} \times \overset{0.000}{9.261} = \overset{0.000}{27.783}$ $-4.5x^2 = -4.5 \times 4.41 = -19.845$ $7.2x = 7.2 \times 2.1 = 15.12$ sum 0.00 0.0 0.00003 0.000 0.000 0.00003 w(2.1) = 44.7306 + 27.783 - 19.845 + 15.12 - 0.1 = 67.6886

Example, cont. I

Let's assume now that coefficients are not exact values. There are "rough values" :

$$a_0 = 2.3, a_1 = 3, a_2 = -4.5, a_3 = 7.2, a_4 = -0.1$$





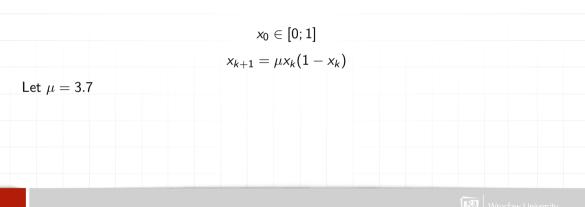
Example, cont. II

Conclusion

When **the input data** are not accurate, increasing the number of significant digits (in calculations) does not increase the accuracy of the result. The rule is known as *Garbage in, garbage out...*



Yet another example





Results

Ν	32 bit	64 bit	128 bit
2	0.2566875	0.2566875	0.2566875
4	0.768053 <mark>4</mark>	0.7680533	0.7680533
6	0.83128 <mark>92</mark>	0.8312889	0.8312889
8	0.923676 <mark>1</mark>	0.9236760	0.9236760
10	0.7133774	0.7133778	0.7133778
12	0.68149 <mark>39</mark>	0.6814953	0.6814953
14	0.58503 <mark>34</mark>	0.5850375	0.5850375
16	0.3381 <mark>789</mark>	0.3381866	0.3381866
18	0.5266 <mark>685</mark>	0.5266460	0.5266460
20	0.2649 <mark>377</mark>	0.2649240	0.2649240
24	0.772 <mark>6893</mark>	0.7727947	0.7727947
28	0.9247 <mark>919</mark>	0.9247575	0.9247575

Ν	32 bit	64 bit	128 bit
32	0.66 <mark>27011</mark>	0.6632869	0.6632869
36	0.26 <mark>65924</mark>	0.2677791	0.2677791
40	0.75 <mark>97211</mark>	0.7502111	0.7502111
45	0.3075923	0.4160662	0.4160662
50	0.8943822	0.9210730	0.9210730
55	0.2570739	0.7404139	0.7404139
60	0.5249998	0.5649204	0.5649208
70	0.9157254	0.6021104	0.6020892
80	0.9222577	0.40 <mark>07202</mark>	0.4019857
90	0.2573895	0.5755109	0.6455021
100	0.7139580	0.3158045	0.8947899
110	0.2567323	0.7575933	0.5199780



Results (computed up to 1000 significant digits)

k	x(k)				
	x(K)				
10:	0.7133778				
	0.264924				
30:	0.7073271				
40:	0.7502111				
50:	0.921073				
60:	0.5649208				
70:	0.6020892				
80:	0.401986				
90:	0.6455194				
100:	0.8950223				
110:	0.5189601				



Computer program

```
#include <stdio.h>
int main()
        float s:
        double d;
        long double e;
        int i:
        s = d = e = 0.5:
        for (i = 1; i <= 110; i ++)
                s = 3.7F * s * (1.F - s);
                d = 3.7 * d * (1. - d);
                 e = 3.7L * e * (1.L - e);
                 if (i%10==0)
                     printf("%10iu%.7fu%.7lfu%.7Lf\n", i, s, d, e);
        return 0:
```



Additional literature I

- Zuber R., Metody numeryczne i programowanie, WSziP 1975, fragmenty: https://kmim.wm.pwr.edu.pl/myszka/wp-content/uploads/sites/2/ 2020/10/zuber.pdf.
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Colophon

Presentation typeset using $\[AT_EX 2_{\mathcal{E}}\]$ system with beamer class using Latin Modern font. Cover page illustration is a part of a picture, showing an 8-bit micro-controller Intel 8742, which contains a single chip CPU with a speed of 12 MHz, 128 bytes of RAM, 2048 bytes EPROM and input/output. Sameli, Ioan. 2006. Old processor. May 25th. Flickr. http://www.flickr.com/photos/biwook/153062645/.

